

Dynamic Programming 2

Problem Solving Club

November 23, 2016

What is dynamic programming?

- Dynamic programming requires recursive thinking
- Wikipedia: “a method for solving a complex problem by breaking it down into a collection of simpler subproblems, solving each of those subproblems just once, and storing their solutions – ideally, using a memory-based data structure”
- Overall, a bit hard to define

Longest Common Subsequence (LCS)

Review from last week's meeting

Given two strings:

$X = \text{bacda}$

$Y = \text{dbdc}$

$LCS(X, Y) = ?$

$LCS\text{-length}(X, Y) = ?$

General steps to solving a DP problem

1. Formulate the problem in terms of a mathematical function
 - Each input corresponds to exactly one output
 - The output of the function depends only on its inputs (no side effects)
 - What would be a function for *LCS-length*?
 - $LCS\text{-length} : (X : \text{string}, Y : \text{string}) \rightarrow \text{integer}$
2. Find a recurrence formula for the problem in terms of smaller subproblem(s)
 - What is the recurrence for *LCS-length*?
 - $LCS\text{-length}(Xa, Ya) = LCS\text{-length}(X, Y) + 1$
 - $LCS\text{-length}(Xa, Yb) = \max[LCS\text{-length}(Xa, Y), LCS\text{-length}(X, Yb)]$
3. Recognize and solve the base cases
 - What are the base cases for *LCS-length*?
 - $LCS\text{-length}(X, \epsilon) = LCS\text{-length}(\epsilon, Y) = 0$
4. Code it

Coding

- Figure how many total states your function has
- This determines how much memory your program will need
- How many states does LCS-length have?
 - $LCS\text{-}length : (X : \text{string}, Y : \text{string}) \rightarrow \text{integer}$
 - $LCS\text{-}length(Xa, Ya) = LCS\text{-}length(X, Y) + 1$
 - $LCS\text{-}length(Xa, Yb) = \max[LCS\text{-}length(Xa, Y), LCS\text{-}length(X, Yb)]$
 - $LCS\text{-}length(X, \epsilon) = LCS\text{-}length(\epsilon, Y) = 0$
- In this particular recurrence, X and Y are always **prefixes** of the original string
- For better runtime performance, define an alternative recurrence:
 - $LCS\text{-}length2 : (x : \text{integer}, y : \text{integer}) \rightarrow \text{integer}$
 - $LCS\text{-}length2(x, y) = LCS\text{-}length2(x-1, y-1) + 1$ if $X[x] = Y[y]$
 - $LCS\text{-}length2(x, y) = \max[LCS\text{-}length2(x, y-1), LCS\text{-}length2(x-1, y)]$ otherwise
 - $LCS\text{-}length2(x, 0) = LCS\text{-}length2(0, y) = 0$
- Is this a mathematical function?

Coding Bottom-up

- For **bottom-up** implementation, you must determine a correct iteration order that processes smaller subproblems before larger ones

- $LCS\text{-}length2 : (x : \text{integer}, y : \text{integer}) \rightarrow \text{integer}$
- $LCS\text{-}length2(x, y) = LCS\text{-}length2(x-1, y-1) + 1$ if $X[x] = Y[y]$
- $LCS\text{-}length2(x, y) = \max[LCS\text{-}length2(x, y-1), LCS\text{-}length2(x-1, y)]$ otherwise
- $LCS\text{-}length2(x, 0) = LCS\text{-}length2(0, y) = 0$

- What would be a correct iteration order for $LCS\text{-}length2$?

- `string X, Y`
- `int dp[|X| + 1][|Y| + 1]`
- `for 0 ≤ x ≤ |X|`
 - `for 0 ≤ y ≤ |Y|`
 - `if x == 0 or y == 0:` `dp[x][y] = 0`
 - `else if X[x] == Y[y]:` `dp[x][y] = dp[x-1][y-1] + 1`
 - `else:` `dp[x][y] = max(dp[x][y-1] + dp[x-1][y])`

- What do we print as the answer?

Coding Top-down

- For **top-down** implementation, it is **not necessary** to find an iteration order
- Allow the computer to do it for you (like Excel, functional programming)
- Implement a function in your program that matches the mathematical function
 - $LCS-length2 : (x : integer, y : integer) \rightarrow integer$
 - $LCS-length2(x, y) = LCS-length2(x-1, y-1) + 1$ if $X[x] = Y[y]$
 $= \max[LCS-length2(x, y-1), LCS-length2(x-1, y)]$ otherwise
 - $LCS-length2(x, 0) = LCS-length2(0, y) = 0$
- `string X, Y`
- `int dp[|X| + 1][|Y| + 1] = initialized to -1`
- `int lcs(int x, int y)`
 - `int ans`
 - `if dp[x][y] != -1: ans = dp[x][y]`
 - `else if x == 0 or y == 0: ans = 0`
 - `else if X[x] == Y[y]: ans = lcs(x-1, y-1) + 1`
 - `else: ans = max(lcs(x, y-1) + lcs(x-1, y))`
 - `dp[x][y] = ans`
 - `return ans`

Coding: Bottom-up vs. Top down

- `string X, Y`
- `int dp[|X| + 1][|Y| + 1]`
- `for 0 ≤ x ≤ |X|`
 - `for 0 ≤ y ≤ |Y|`
 - `if x == 0 or y == 0:` `dp[x][y] = 0`
 - `else if X[x] == Y[y]:` `dp[x][y] = dp[x-1][y-1] + 1`
 - `else:` `dp[x][y] = max(dp[x][y-1] + dp[x-1][y])`

- `string X, Y`
- `int dp[|X| + 1][|Y| + 1] = initialized to -1`
- `int lcs(int x, int y)`
 - `int ans`
 - `if dp[x][y] != -1:` `ans = dp[x][y]`
 - `else if x == 0 or y == 0:` `ans = 0`
 - `else if X[x] == Y[y]:` `ans = lcs(x-1, y-1) + 1`
 - `else:` `ans = max(lcs(x, y-1) + lcs(x-1, y))`
 - `dp[x][y] = ans`
 - `return ans`

Coding: Bottom-up vs. Top down

What might be some reasons to prefer one method over the other?

- Runtime performance
 - A complicated issue
 - Bottom-up computes **all states**, while top-down only computes **relevant states**
 - If most states are visited, top-down is usually slower than bottom-up due to call stack
 - Top-down approach can cause a stack overflow
- Ease of coding
 - Bottom-up requires determination of the iteration order
 - For some types of problems (e.g. travelling salesman), the iteration order is non-obvious
- Personal preference